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#### WINTER-18 EXAMINATION

Subject Code: 17330 **Subject Name: Data Structure Model Answer** 

### **Important Instructions to examiners:**

- 1) The answers should be examined by key words and not as word-to-word as given in the model answer scheme.
- 2) The model answer and the answer written by candidate may vary but the examiner may try to assess the understanding level of the candidate.
- 3) The language errors such as grammatical, spelling errors should not be given more Importance (Not applicable for subject English and Communication Skills.
- 4) While assessing figures, examiner may give credit for principal components indicated in the figure. The figures drawn by candidate and model answer may vary. The examiner may give credit for any equivalent figure drawn.
- 5) Credits may be given step wise for numerical problems. In some cases, the assumed constant values may vary and there may be some difference in the candidate's answers and model answer.
- 6) In case of some questions credit may be given by judgement on part of examiner of relevant answer based on candidate's understanding.
- 7) For programming language papers, credit may be given to any other program based on equivalent concept.

Q. No.	Sub Q. N.	Answer	Marking Scheme
1	A	Attempt any SIX :	12 M
	a	Define data structure. Enlist any two operations on it.	2 M
	Ans	<b>Data structure:</b> A Data structure is logical and mathematical model used for storing and organizing data in particular way so that it can be accessed efficiently.	Definition: 1M, any two operations: 1M
		OR	
		<b>Data structure</b> is a collection of data values, the relationships among them, and the functions or operations that can be applied to the data.	
		Operations perform on data structure:	
		1. Insertion	
		2. Deletion	
		3. Searching	
		4. Sorting	
		5. Traversing	
		6. Merging	



b	Define sorting. Write its types.	2 M
Ans	<b>Sorting</b> is the process of placing elements from a collection in specific logical order. For example, a list of words could be sorted alphabetically or by length.	Definition: 1M, any two types: 1M
	Types:	
	1. Bubble sort	
	2. Selection sort	
	3. Insertion sort	
	4. Quick sort	
	5. Radix sort	
	6. Merge sort	
	7. Shell sort	
c	Write any two applications of stack.	2 M
Ans	1. Reversing a list	Each
	2. Conversion of infix expression to postfix expression	application: 1M
	3. Evaluation of postfix expression	
	4. Conversion of infix expression to prefix expression	
	5. Evaluation of prefix expression	
	6. Recursion	
d	Write any four primitive operations on queue.	2 M
Ans	1. <b>enqueue</b> – add (store) an item to the queue from rear end	Listing each
	2. <b>dequeue</b> – remove an item from the queue from front end.	operation: ½ M
	3. <b>peek</b> – Gets the element at the front of the queue without removing it.	
	4. <b>isfull</b> – Checks if the queue is full.	
	5. <b>isempty</b> – Checks if the queue is empty.	
e	Define the terms NULL pointer and next pointer for linked list.	2 M
Ans	<b>NULL pointer:</b> It is an address field of the last node in the linked list which specifies end of the list by holding value as a NULL.	Each 1M
	Next pointer: It is one of field of node structure which contains an address of	



	the next node.	
f	Define binary search tree.	2 M
Ans	A Binary Search Tree (BST) is a tree in which all the nodes follow the belowmentioned properties –	Correct definition: 2M
	• The left sub-tree of a node has a key less than or equal to its parent node's key.	
	• The right sub-tree of a node has a key greater than to its parent node's key.	
g	Write any two applications of graph.	2 M
Ans	Applications of graphs:	Each 1M
	1. To represent road map	
	2. To represent circuit or networks	
	3. To represent program flow analysis	
	4. To represent transport network	
	5. To represent social network	
	6. Neural networks	
h	Define the term recursion.	2 M
Ans	Recursion is the process of calling function by itself till terminating condition.	Correct definition: 2M
В	Attempt any TWO:	8 M
a	Define the following terms with respect to tree:	4 M
	i) Leaf node	
	ii) Degree of node	
	iii) Height of tree	
	iv) Descendant node	
Ans	i. <b>Leaf node:</b> A node having no child or of degree zero is called a terminal node or leaf node.	Each 1 M
	ii. <b>Degree of node</b> : It is defined as maximum number of child nodes of any node.	
	OR	
	Degree of node is the number of nodes connected to a particular node.	



	iii. <b>Height of tree:</b> The longest path from root node to the terminal node is known as Depth of tree or height of tree.	
	iv. <b>Descendant node:</b> A descendant refers to any element that is connected lower down the hierarchy tree – no matter how many levels lower. A descendant is a child, grandchild, great-grandchild, and so on.	
b	Write a 'program in c' language for selection sort.	4 M
Ans	#include <stdio.h></stdio.h>	Correct Logic
	#include <conio.h></conio.h>	2M,Correct Syntax 2M
	void main()	**any other
	{	relevant logic can be
	int a[100],n,i,j,min,temp;	considered**
	clrscr();	
	printf("\n Enter the Number of Elements: ");	
	scanf("%d",&n);	
	printf("\n Enter %d Elements: ",n);	
	for(i=0;i <n;i++)< td=""><td></td></n;i++)<>	
	{	
	scanf("%d",&a[i]);	
	}	
	for(i=0;i< n-1;i++)	
	{	
	min=i;	
	for(j=i+1;j< n;j++)	
	{	
	if(a[min]>a[j])	
	min=j;	
	}	
	if(min!=i)	
	{	



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```
temp=a[i];
                     a[i]=a[min];
                     a[min]=temp;
                   }
                }
                printf("\n The Sorted array in ascending order: ");
                for(i=0;i<n;i++)
                  printf("%d ",a[i]);
                getch();
              }
      Convert the given infix expression into postfix using stack and write down step of
                                                                                                 4 M
c
      conversion
               a \uparrow b * c - d + e.
                                                                                                 Correct
Ans
                                                                                                  answer: 4M,
       Symbol Scanned
                                     Stack
                                                                  Expression
                                                                                                 *give
                                                                                                 appropriate
                                                                                                 marks for steps
       а
                                                                  а
        \uparrow
                                     (个
                                     (个
       b
                                                                  ab
                                     (*
                                                                  ab个
                                     (*
                                                                  ab↑c
       С
                                                                  ab个c*
                                     (-
                                                                  ab↑c*d
                                     (-
                                                                  ab↑c*d-
                                     (+
                                                                  ab↑c*d-e
       e
                                                                  ab↑c*d-e+
                                     empty
```



		Ans: a b ↑ c * d – e +	
2		Attempt any FOUR:	16 M
	a	Define Algorithm. Describe different approaches for designing an algorithm.	4 M
	Ans	Algorithm: It is sequence of steps/instructions that must be followed to solve a problem. In other words, an algorithm is a logical representation of the steps/instructions which should be executed to perform a meaningful task.  Approaches to design an algorithm:  1. Top-down approach	Definition: 1M Diagram: 1M, Explanation: 2M
		2. Bottom-up approach	
		Top-down approach:	
		a. A top-down design approach starts by dividing complex algorithm into one or more modules or subsystems	
		b. Each subsystem is then refined in yet greater detail, sometimes in many additional sub system levels, until the entire specification is reduced to base elements. c. Top-down design method is a form stepwise refinement where we begin with the Top most modules and incrementally add modules that it calls.	
		2. Bottom-up approach:	
		a. In this approach the individual base elements of the system are first specified in detail.	
		b. These elements are then linked together to form larger subsystems, which then in turn are clubbed in many levels, until a complete top-level system is formed.	
		Top-down approach  Module 1  Module 2  Module n  Each module can be divided into one or more sub modules	



b	Write difference between stack and queue (any 4 points).				
Ans	Stack  1. In Stack insertion and deletion operations are performed at same end.	Queue  1. In Queue insertion and deletion operations are performed at different end.	Correct poin 1 M		
	2. In stack the element which is inserted last is first to delete so it is called Last In First Out.	2. In Queue the element which is inserted first is first to delete so it is called First In First Out.			
	<b>3.</b> In stack only one pointer is used called as Top.	<b>3.</b> In Queue two pointers are used called as front and rear.			
	4. In Stack Memory is not wasted	4. In Queue memory can be wasted/ unusable in case of linear queue.			
	5. Stack of books is an example of stack	<b>5.</b> Students standing in a line at fees counter is an example of queue			
	<ul><li>6. Application:</li><li>Recursion</li><li>Polish notation</li></ul>	<ul><li>6. Application:</li><li>In computer system for organizing processes.</li></ul>			
		<ul> <li>In mobile device for sending and receiving messages</li> </ul>			
	Write an algorithm to POP an element from stack.				
С	write an algorithm to 1 Or an element	from stack.	4 M		
Ans	POP: The process of deleting an element POP operation. As deletion takes place	from top of the stack is called from the top of the stack. After ever	Correct		
	POP: The process of deleting an element POP operation. As deletion takes place operation the top of the stack is decrement	from top of the stack is called from the top of the stack. After ever	Correct		
	POP: The process of deleting an element POP operation. As deletion takes place operation the top of the stack is decrement population.	from top of the stack is called from the top of the stack. After every sted by one.  Operation  Top of the stack is called from the top of the stack. After every sted by one.  Top of the stack is called from the stack. After every sted by one.	Correct algorithm: 4		
	POP: The process of deleting an element POP operation. As deletion takes place operation the top of the stack is decrement population.  **Top	from top of the stack is called from the top of the stack. After every sted by one.  Operation  Top of the stack is called from the top of the stack. After every sted by one.  Top of the stack is called from the stack. After every sted by one.	Correct algorithm: 4		
	POP: The process of deleting an element POP operation. As deletion takes place operation the top of the stack is decrement population.  Algorithm for pop operation:  First check for underflow i.e. if top is -1 is operation cannot be performed. If there is removes an element placed at top of the stack is decrement. If TOP = =-1, then	from top of the stack is called from the top of the stack. After every sted by one.  Operation  Top of the stack is called from the top of the stack. After every sted by one.  Top of the stack is called from the stack. After every sted by one.	Correct algorithm: 4		

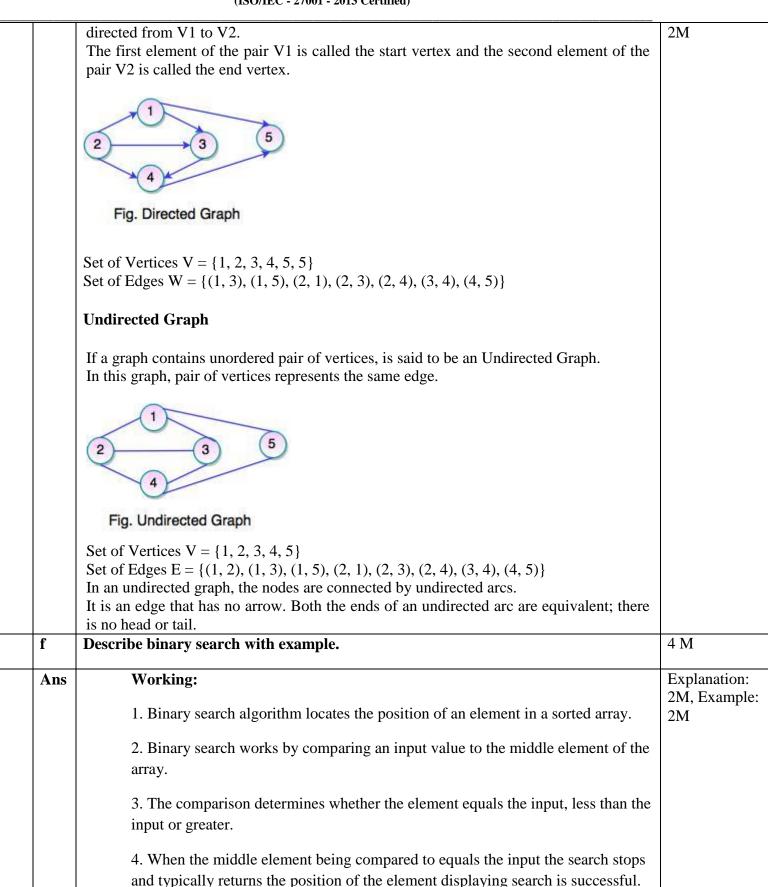


	5.  TOP = TOP - 1	
	6. Exit.	
d	Create a binary search tree for the following data:	4 M
	10, 25, 15, 5, 2, 7, 12	
Ans	Step 1:  Step 6:  Step 2:  Step 3:  Step 4:  Step 5:	Correct answer: 4M, give marks to right steps
e	Describe directed and undirected graph with suitable example.	4 M
Ans	Directed Graph	Directed Graph Explanation: 2M
	If a graph contains ordered pair of vertices, is said to be a Directed Graph.	Undirected
	If an edge is represented using a pair of vertices (V1, V2), the edge is said to be	Graph Explanation:



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5. If the middle element is not equal to the input then a comparison is made to



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determine whether the input is less than or greater than the middle element.

- 6. Accordingly given input array is divided into two subarrays, lower subarray containing elements less than the middle element and upper subarray containing elements greater than the middle element.
- 7. This process continues from step 2-6 till either the input value is found or the search is unsuccessful.

#### **Searching Element in Given List:**

List: 23, 12, 5, 29, 10, 65, 55, 70

Pre-condition for Binary search is array elements must be in ascending order.

The given list is not sorted.

Sorted List A= {5, 10, 12, 23, 29, 55, 65, 70}

Search element (k) = 65

i. Low=0 high=7

Mid=(0+7)/2=3

A[mid] = a[3] = 23

65>23

k>a [mid]

ii. Low=mid+1 High=7 Mid= (4+7)/2=5

A[mid] = a [5] = 29

65>29

iii. Low=6 high=7

Mid = (6+7)/2 = 6

A[mid] = a [6] = 65

A[mid] = k

Therefore key element is found at 6th position, no. of comparison required = 3.



	Search is successful			
	Attempt any FOUR :			16 M
a	Describe the concept of time complex	kity with exa	nple.	4 M
Ans	Time Complexity: Time complexity of time that it needs to run to completion. concentrate on developing only frequent Example: #include <stdio.h> void main () {   int i, n, sum, x;   sum=0;   printf("\n Enter no of data to be added" scanf("% d", &amp;n);   for(i=0; i<n; ",="" &x);="" calculation="" computational="" i++)="" of="" printf("\n="" scanf("%d",="" sum="%d" sum);="" td="" time:<="" {="" }=""><td></td></n;></stdio.h>			
	Statement	Frequenc	Computational Time	
	sum=0	1 1	t <sub>1</sub>	
	printf("\n Enter no of data to be added")	1	t <sub>2</sub>	
	scanf("% d", &n)	1	t <sub>3</sub>	
	for(i=0; i <n; i++)<="" td=""><td>n+1</td><td>(n+1)t<sub>4</sub></td><td></td></n;>	n+1	(n+1)t <sub>4</sub>	
	scanf("%d", &x)	n	nt <sub>5</sub>	
	sum=sum+x	n	nt <sub>6</sub>	
	printf("\n Sum = %d ", sum)	1	t <sub>7</sub>	
	Total computational time= $t1+t2+t3+(n+t)$ T = n(t4+t5+t6)+(t1+t2+t3+t4+t7) For large n, T can be approximated to $T = n(t4+t5+t6) = kn$ where $k = t4+t5+t6$ Thus $T = kn$ or $T$	ŕ	ut5+t7	
b	Write an algorithm for traversal of g	od. 4 M		
ש			, <b>-</b>	



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The aim of DFS algorithm is to traverse the graph in such a way that it tries to go far from the root node. Stack is used in the implementation of the depth first search. Back tracking used in this algorithm.

Algorithm: 2 marks

Example 2Marks

#### Algorithm

Step1: Start

Step2: Initialize all nodes as unvisited

Step3: Push the starting node onto the stack. Mark it as waiting.

Step4: Pop the top node from stack and mark is as visited. Push all its adjacent nodes

into the stack &mark them as waiting.

Step 5 .Repeat step 4 until stack is empty.

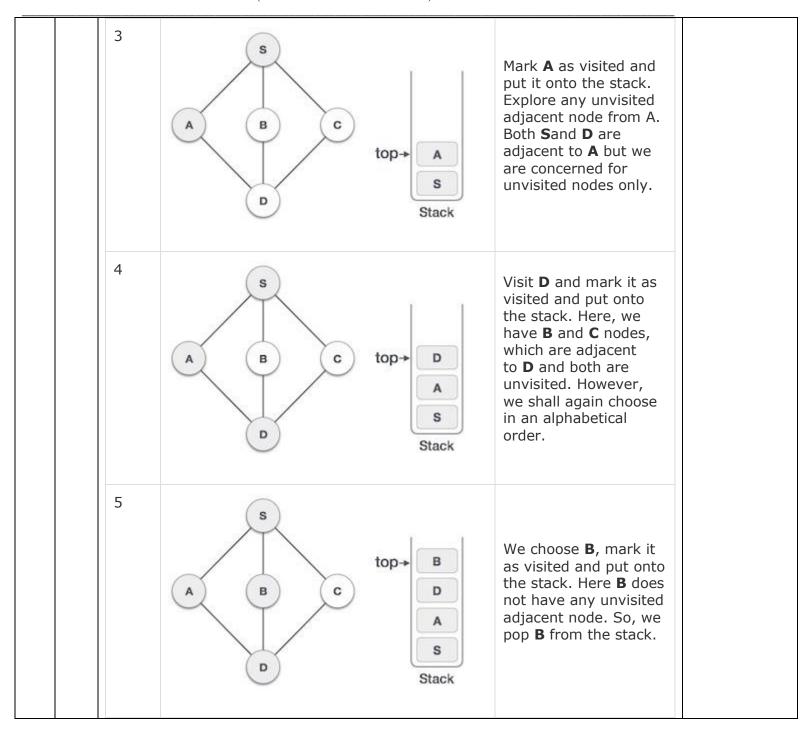
Step 6: Stop

Step	Traversal	Description
1	A B C Stack	Initialize the stack.
2	A B C top→S Stack	Mark <b>S</b> as visited and put it onto the stack. Explore any unvisited adjacent node from <b>S</b> . We have three nodes and we can pick any of them. For this example, we shall take the node in an alphabetical order.



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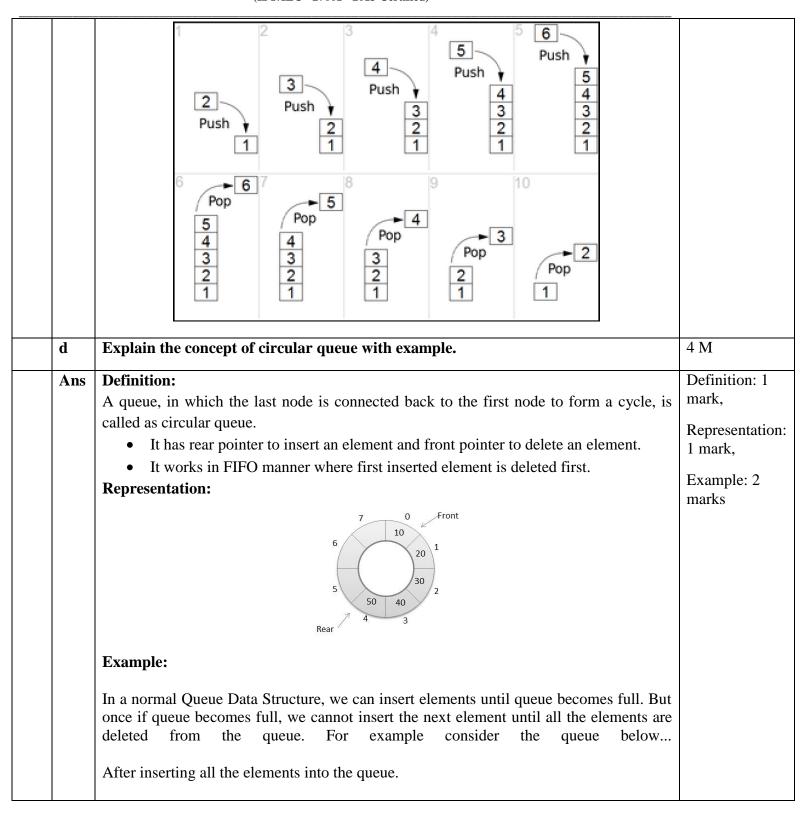




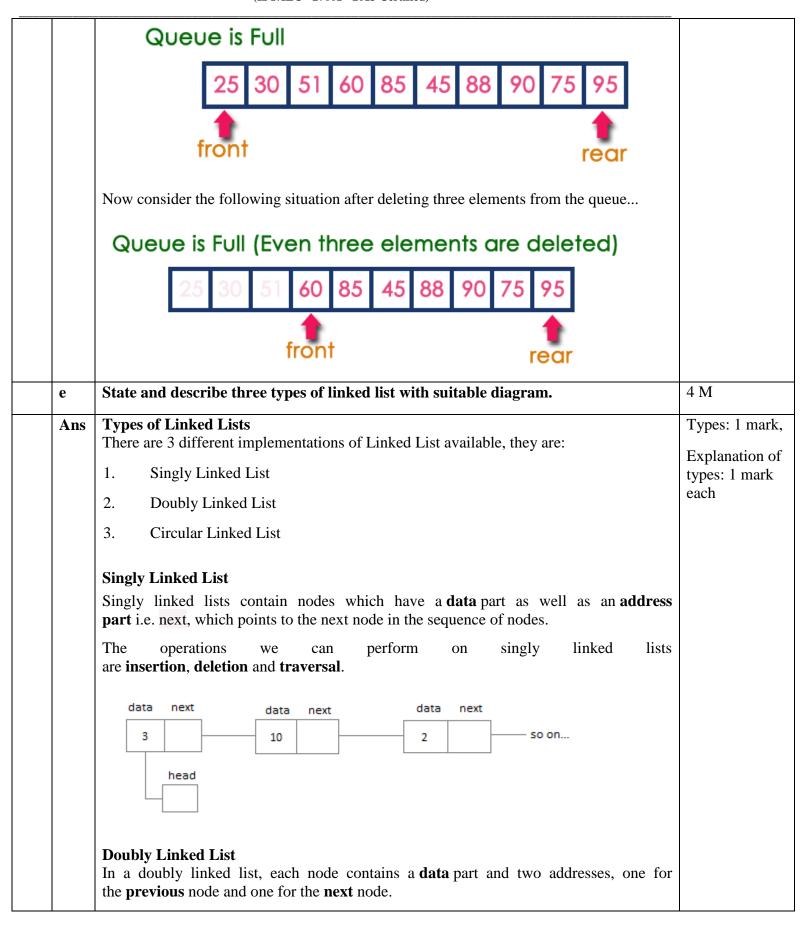
	We check the stack top for return to the previous node and check if it has any unvisited nodes. Here, we find <b>D</b> to be on the top of the stack.	
	Only unvisited adjacent node is from <b>D</b> is <b>C</b> now. So we visit <b>C</b> , mark it as visited and put it onto the stack.	
c	Describe with example, use of stack in reversing a list.	4 M
Ans	Reversing a list  To reverse a list, the elements of list are pushed onto the stack one by one. Once all elements are pushed on the stack they are popped one by one. Since the element last pushed in comes out first, hence reversal of string occurs.  Example: a list contains elements as {1, 2, 3, 4, 5, 6}. Every push operator will push an element on top of stack. Once all elements are pushed one can pop all elements and save it which results in to reversing of list as {6, 5, 4, 3, 2, 1}.	Explanation : 2 marks, Example: 2 marks



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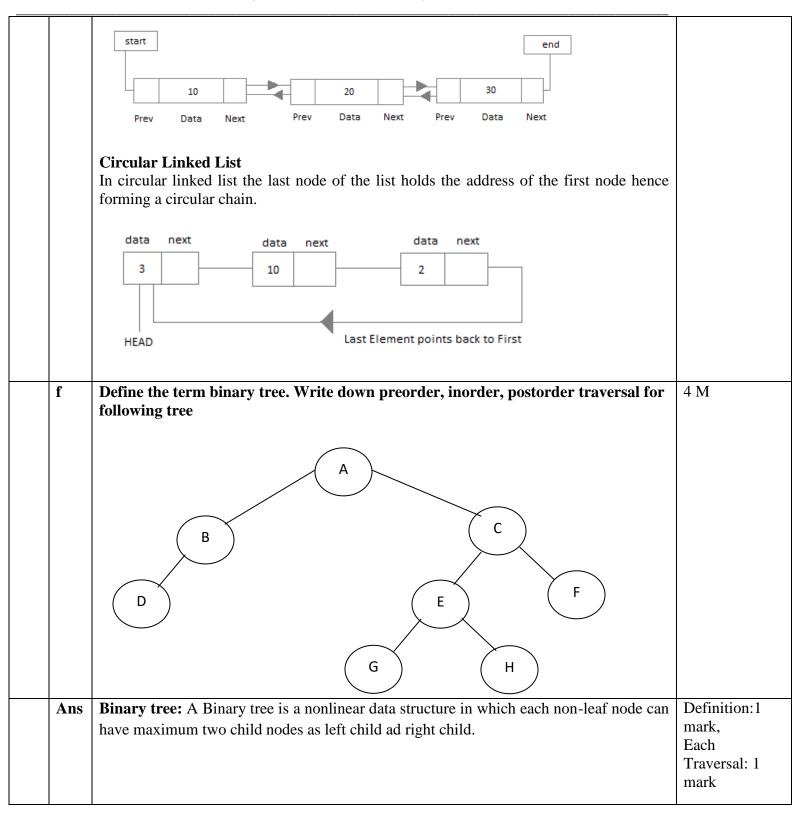


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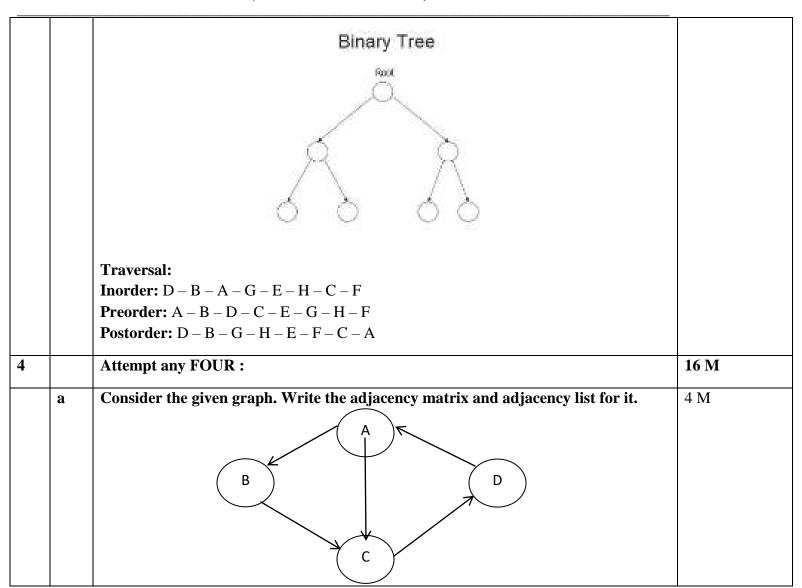


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Ans	Adjacency matrix	A for the above gra	ıph			Adjacency Matrix: 2
		A	В	C	D	Marks, Adjacency list:
	A	0	1	1	0	2 Marks
	В	0	0	1	0	
	C	0	0	0	1	
	D	1	0	0	0	
	B C D Adjacency List	B   C   D   A		C		
		Node	Adjacenc	y List		
		A	В,С			
		В	С			
		С	D			
		D	A			
b		een general tree and	•			4 M



	Sr.no	General Tree	Binary Tree	Any four
	1	A general tree is a data	A Binary tree is a data	correct
		structure in that each node		points: 1 mark
		can have infinite number of		each
	2	children In general tree, root has in-	and right	
	2	degree 0 and maximum out-	In binary tree, root has in- degree 0 and maximum out-	
		degree n.	degree 2.	
	3	In general tree, each <b>node</b>	In binary tree, each node	
		have in-degree one and	have in-degree one and	
		maximum out-degree n.	maximum out-degree 2.	
	4	Height of a general tree is	Height of a binary tree is :	
		the length of longest path		
		from root to the leaf of tree.	(Height(Left Child) ,	
		Height(T) =	Height(Right Child) + 1}	
		{max(height(child1) ,		
		height(child n) +1)		
	5	height(child-n))+1} Subtree of general tree are	Subtree of binary tree is	
	3	not ordered	ordered.	
		General tree	SEAL SECTION OF THE SEAL S	
		501101411100	Binary Tree	
		R	R	
		999999		
c	Explain t	he procedure for deleting first noo	de from a singly linked list.	4 M
c Ans		he procedure for deleting first nod	le from a singly linked list.	4 M 4 marks for
	Deleting f			
	Deleting f  We can us  Ste Ste po Ste wii Ste Ste Ste	First node from singly linked list  the the following steps to delete a first  the p 1: Check whether list is Empty (  the p 2: If it is Empty then, displet singleter and terminate the function.  The p 3: If it is Not Empty then, deep 4: Check whether list is having of the p 5: If it is	t node from singly linked list  head == NULL)  ay 'List is Empty!!! Deletion is not  fine a Node pointer 'temp' and initialize  nly one node (temp -> next == NULL)  TRUE then set head = NULL and	4 marks for correct
	Deleting f  We can us  Sto Sto yo Sto Sto Wi Sto Sto George	First node from singly linked list  the the following steps to delete a first  the p 1: Check whether list is Empty (  the p 2: If it is Empty then, displaysible and terminate the function.  The p 3: If it is Not Empty then, deep 4: Check whether list is having of the p 5: If it is lette temp (Setting Empty list conditions).	t node from singly linked list  head == NULL)  ay 'List is Empty!!! Deletion is not  fine a Node pointer 'temp' and initialize  nly one node (temp \rightarrow next == NULL)  TRUE then set head = NULL and tions)	4 marks for correct
	Deleting f  We can us  Sto Sto yo Sto Sto Wi Sto Sto George	First node from singly linked list  the the following steps to delete a first  the p 1: Check whether list is Empty (  the p 2: If it is Empty then, displet singleter and terminate the function.  The p 3: If it is Not Empty then, deep 4: Check whether list is having of the p 5: If it is	t node from singly linked list  head == NULL)  ay 'List is Empty!!! Deletion is not  fine a Node pointer 'temp' and initialize  nly one node (temp \rightarrow next == NULL)  TRUE then set head = NULL and tions)	4 marks for correct



Ans	A priority Queue is a collection of elements where each element is assigned a priority and the order in which elements are added into the queue.	Explanation :2M
	The rules for processing the elements of priority queue are: 1) An element with higher priority is processed before any element of lower priority. 2) Two elements with the same priority are processed according to the order in which they are added to the queue (FCFS).	Example :2M
	One of the examples of priority queue is a queue used in operating system.	
	The operating system has to handle a large number of jobs. These jobs have to be properly scheduled. The operating system assigns priorities to each type of job. The jobs are placed in a queue and the job with the highest priority will be executed first.	
	Example:	
	Array representation:	
	Array element of priority queue has a structure with data, priority and order.	
	Priority queue with 5 elements is as shown below:-	
	C,1,4 B,3,2 B,3,5 A,4,1 D,5,3	
	In the above diagram, each structure element has three members as information, priority and order in which element is arrived in the list.	
e	Perform bubble sort on following data to sort all elements in ascending order.	4 M
	15, 10, 02, 35, 08 (show all steps).	



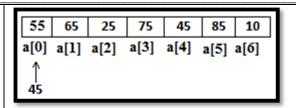
Ans	Pass 1: 15 10 02 35 08	Correct answer 4 M
	10 15 02 35 08 10 02 15 35 08 10 02 15 35 08	** Give Stepwise Marks**
	Pass 2:  10 02 15 08 35  10 02 15 08 35  02 10 15 08 35  02 10 08 15 35  Pass3: 02 10 08 15 35  02 10 08 15 35	
	02 08 10 15 35  Pass 4:  02 08 10 15 35  The Final sorted array in ascending order is  02,08,10,15,35.	
f	Write a C program to calculate the factorial of a number using recursion.	4 M
Ans	#include <stdio.h> #include<conio.h></conio.h></stdio.h>	syntax: 2 marks, Logic: 2 marks



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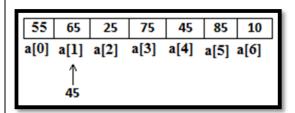
```
int Fibonacci(int);
      int main()
      {
      int n, i = 0, c;
      scanf("%d",&n);
      printf("Fibonacci series\n");
      for (c = 1; c \le n; c++)
      printf("%d\n", Fibonacci(i));
      i++;
      getch();
      return 0;
      int Fibonacci(int n)
      if (n == 0)
      return 0;
      else if (n == 1)
      return 1;
      else
      return ( Fibonacci(n-1) + Fibonacci(n-2) );
      Attempt any FOUR:
                                                                                              16 M
      Consider the following array:
                                                                                              4 M
a
      55 65 25 75 45 85 10
      Write stepwise procedure to find 45 using linear search.
      Step 1: Compare element 45 with a [0].
                                                                                              Correct answer
Ans
                                                                                              4M
```

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Since, 45 is not equal to a[0], compare next element.

**Step 2:** Compare element 45 with a[1].



Since, 45 is not equal to a[1], compare next element.

**Step 3:** Compare element 45 with a[2].

55	65	25	75	45	85	10
a[0]	a[1]	a[2]  ↑ 45	a[3]	a[4]	a[5]	a[6]

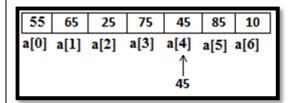
Since, 45 is not equal to a[2], compare next element.

**Step 4:** Compare element 45 with a[3].

55	65	25	75	45	85	10
a[0]	a[1]	a[2]	a[3]  ↑ 45	a[4]	a[5]	a[6]

Since, 45 is not equal to a[3], compare next element.

**Step 4:** Compare element 45 with a[4].



Since, 45 is equal to a[4]. Therefore we conclude that key element 45 is found.



b	Define linked list. Write its two advantages and disadvantages.	4 M
Ans	<b>Linked list:</b> It is linear collection of data elements. Each element in linked list is called as 'node'. Each node contains two fields. First is INFO which stores data & second is NEXT which is linked with address of next node in a list.	Define 2M, Advantages 1M and
	START NODE 1 NODE 2 NODE 3  10 20 30 NULL INFO NEXT INFO NEXT	Disadvantages 1M)
	Advantages:	
	<ol> <li>Linked list is a dynamic memory allocation</li> <li>In Linked list Insertions and deletions can be done easily just by storing an address of new node into previously existing node ie without movement of elements for insertion and deletion.</li> <li>Disadvantages:</li> </ol>	
	<ol> <li>It requires more space as pointers are also stored along with information.</li> <li>Different amount of time is required to access each element.</li> <li>If we have to go to a particular element then we have to go through all those elements that come before that element.</li> <li>We cannot traverse it from last and only from the beginning.</li> <li>It is not easy to sort the elements stored in the linear linked list.</li> </ol>	
c	Write algorithm for inorder traversal of binary tree.	4 M
Ans	The algorithm works by:  1. Traversing the left sub-tree	4M for correct algorithm
	2. Visiting the root node, and finally	
	3. Traversing the right sub-tree.	
	OR	
	INORDER(ROOT)	
	Step 1: Repeat Steps 2 to 4 while (ROOT != NULL)	
	Step 2: INORDER(ROOT-> LEFT)	
	Step 3: Write ROOT->DATA	



	Step 4: 1	NORDF		RIGHT) [END				
	Step 5: 1		21(110 0 1 7	1110111) [2112	01 20 01 ]			
d	1		<sup>2</sup> notation	Also give exam	nlo			4 M
						.1 C	1 '4	
Ans	of an alg		is used in C	computer Science	ce to describe	the performan	nce or complexity	2M for explanation
	2 Rig O	specifics	ally describe	es the worst-case	e scenario and	d can be used	to describe the	and 2M for correct
	_	-	•	ne space used (e.				example
	3.Exam	ple:						
	Ş	Since Big	g O notation	is a f(n) where	n = no. of inp	uts		
	(	Consider	$f(n) = 10n^2 +$	-4n+2, $g(n) = 1$	$n^2$ c=11 wher	e c = constant	:>0 and n>1	
	Г	n	1	2	3	4	5	
		f(n)	16	50	104	178	272	
		g(n)	1	4	9	16	25	
		c. g(n)	11	44	99	176	275	
	1 1							
e	Define l	Hash fur	nction. Expl	ain any one me	ethod of hash	ing.		4 M
Ans	C	data of a	fixed size. T	ny function that The values return digests, or simp	ned by a hash		arbitrary size to	2M definition 2M for explanation o
				ten used in com puter software f			a common data	method.
			ctions accele n a large file	erate table or da	tabase lookup	by detecting	duplicated	
	4. I	_	methods:-					
			ivision meth					
		D. IVI	lid square m	emou				

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c. Folding method

5. **Division method:** In this method hash address is calculated by dividing the key value by a prime number or a number without small divisor.

Formula:  $H(K)=K \pmod m$  or  $H(K)=K \pmod m+1$  K- Specify unique key value. m- Specify is a prime number or number without small divisors. Example:- $H(3205)=3205 \mod 97=42$ .

- 6. **Middle square method:** In this method hash address is calculated by taking two digits from middle of square of key value.
- 7. Formula: H(K)=I I-specify digits after deleting digits from both ends of K2 Example:- H(3205)=(3205) 2 =10272015=72
- 8. **Folding method:** In this method hash address is calculated by partitioning key into multiple parts and performing addition. Formula: H(K)=K1+K2+...Kn Example: H(3205)=32+05=37

Given a hash table of 100 locations, calculate the hash value using folding method for keys 5678, 321, and 34567

#### **Solution**

Since there are 100 memory locations to address, we will break the key into parts where each part (except the last) will contain two digits. The hash values can be obtained as shown below:

Key	5678	321	34567
Parts	56 and 78	32 and 1	34, 56 and 7
Sum	134	33	97
Hash Value	34 (Ignore the last carry)	33	97

#### f Describe stack as ADT.

4 M

**Ans** The stack, an abstract data type is defined as an ordered collection of items where items are added to and removed from the end called the "top."

4M for correct explanation

Stacks are ordered LIFO.

The stack operations are given below.

Stack () creates a new stack that is empty. It needs no parameters and returns an empty stack.

push (item) adds a new item to the top of the stack. It needs the item and returns nothing.

Pop () removes the top item from the stack. It needs no parameters and returns the item. The stack is modified.

Peek () returns the top item from the stack but does not remove it. It needs no parameters. The stack is not modified.

isEmpty() tests to see whether the stack is empty. It needs no parameters and returns a boolean value.

size() returns the number of items on the stack. It needs no parameters and returns an integer.



	Attempt	any FO	UR:											16 M
a	Describe	e workin	ng of r	radix	sort v	with e	xample	·•						4 M
Ans		n this me				, ,					-		All the	2M for description 2M for
	d	n pass 1, igit. Afte th bucke	er plac			_					_	_		example
	e	n pass 2, ach elem n bucket.	nent. I		_				_		_	_		
	tl	at the end ne next p umber o	oass. T	Γotal	numb	er of	passes 1	equire	d for s	orting	is equa		-	
	5. L	ast pass	give	s soi	rted lis	st afte	er readi	ng all	eleme	nts fro	m Oth	bucket	to 9th	
		ucket.												
	b	ucket. Example:	18,25	53, 10	000,2,	80,75	,58							
	b	Example:						ng to ui	nit plac	e.				
	6. E	Example:						ng to un	nit plac	ee. 8	9	]		
	6. E	Example:	s arra	nge t	he ele	ment a	accordii	1	nit plac		9			
	6. E Pass1: In	Example:	s arra	nge t	the elements and the second se	ment a	accordii	1	nit plac	8	9			
	6. E Pass1: In  18  253	Example:  this pas	s arra	nge t	he ele	ment a	accordii	1	7	8	9			
	6. E Pass1: In  18  253  1000	Example:	s arra	nge t	the elements and the second se	ment a	accordii	1	7	8	9			
	6. E Pass1: In  18  253  1000  2	this pas	s arra	nge t	the elements and the second se	ment a	accordii	1	7	8	9			
	6. E Pass1: In  18  253  1000  2  80	Example:  this pas	s arra	nge t	the elements and the second se	ment a	5	1	7	8	9			
	6. E Pass1: In  18  253  1000  2	this pas	s arra	nge t	the elements and the second se	ment a	accordii	1	7	8	9			



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	0	1	2	3	4	5	6	7	8	9
1000	1000									
80									80	
2	2									
253						253				
75								75		
18		18								
58						58				

Output of 2nd pass: 1000, 2, 18,253, 58, 75, 80

Pass 3: In this pass arrange the element according to hundred's place.

	0	1	2	3	4	5	6	7	8	9
1000	1000									
2	2									
18	18									
253			253							
58	58									
75	75									
80	80									

Output of 3rd pass: 1000, 2, 18, 58, 75, 80,253

Pass 4: In this pass arrange the element according to thousand's place



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1		0	1	2	3	4	5	6	7	8	9		
	1000		1000										
	2	2											
	18	18											
	58	58											
	75	75											
	80	80				<u> </u>							
	253												
	253	253											
	Output o	of 4th pa	ass: 2,18	3,58,75	5,80,25	53,1000	)						
	Elements	-						.1000					
b	Explain								le diso	ram			4 M
Ans											nto i o	stack is full	2M for stack
AllS												in overflow	overflow and
	and pusi	i opera	tion is c										
	state.	i opera	tion is c										
		i opera	tion is c										2M for stack underflow
												_	
			1 = 4					Mo	0X = 4	1-		7	
	state.				v			Ma	0X = 4	1			
	state.	Max		Slac		usų	E	3	D X = 4	1	stack		
	state.	Max		star		USH	E :	1		1	stact top		
	state.	Max D C B A	= A	star		USH	,	3 2	D C B	-	top		
	state.	Max D C B A		star		USH	,	3	D C B	-	top		
	state.	Max D C B A	- A full	Star top	, =	$\Rightarrow$	,	Star	D C B A	enf	top	pty and pop	



	Max = 4  Pop  Stack empty  Max = 4  Pop  Stack top  Stack empty  underflow	
С	Write an algorithm for traversing in linked list.	4 M
Ans	Step 1: [INITIALIZE] SET PTR = START  Step 2: Repeat Steps 3 and 4 while PTR != NULL  Step 3: Apply Process to PTR DATA	4M for correct algorithm
	Step 4: SET PTR = PTR NEXT [END OF LOOP] Step 5: EXIT	
d	Explain Double Ended Queue with suitable diagram.	4 M
Ans	1. A double-ended queue or dequeue is an abstract data structure that implements a queue for which elements can only be added to or removed from the front (head) or back (tail).	2M for explanation and 2M for diagram
	2. It is also often called a head-tail linked list.	
	3. Dequeue is a special type of data structure in which insertions and deletions will be done either at the front end or at the rear end of the queue.	
	4. The operations that can be performed on dequeues are:	
	a. Insert an item from front end	
	b. Insert an item from rear end	
	c. Delete an item from front end	
	d. Delete an item from rear end	



	Rear Front Front Rear	
e	Draw tree structure for following expression: $(2x + y^2 + z^3) + (3a + 4b + c^2)$ .	4 M
Ans		4M for correstructure
f	Define the following terms with respect to graph:	4 M
	a) Successor	
	b) Indegree	
	c) Path	
	d) Weighted graph.	
Ans	a) Successor: Successor is a node which comes after a particular node. In the below given diagram, successor of node X is V and Y.	Definition ea 1M
	b) <b>Indegree:</b> It is number of edges coming towards a specified node i.e. number of edges that have that specified node as the head. The Indegree of node Y is 3.	
	c) <b>Path:</b> Path is a sequence of alternating vertices and edges such that each successive vertex is connected by the edge. Frequently only the vertices are listed especially if there are no parallel edges. Suppose if we want to find the path to $Z$ from $X$ . The solution is $X-Y-W-Z$ and another path is $X-V-Z$ .	
	d) Weighted graph: A graph whose edges are assigned non-negative numerical values	



